

# PS Algorithms and Data Structures 2026

## Task sheet 7

### Task 19

Consider a hash table of size  $m$  using chaining for collision resolution. After inserting  $n$  elements, the expected length of each chain is  $\frac{n}{m}$ . This assumes simple uniform hashing.

Show that: For any constant  $c > 0$ , the probability that a chain contains at least  $c \cdot \frac{n}{m}$  elements is at most  $\frac{1}{c}$ . In doing so, you apply Markov's inequality and prove it.

### Task 20

Let  $G$  be a directed graph represented in (a) adjacency list format and (b) adjacency matrix format. For each of the following problems, determine the tight runtime complexity in  $\Theta$ -notation based on the following parameters: (1) number of vertices  $n$ , (2) number of edges  $m$ , (3) out-degree of a vertex  $\deg^+(v)$ , (4) maximum out-degree  $\Delta^+$ , (5) in-degree of a vertex  $\deg^-(v)$ , and (6) maximum in-degree  $\Delta^-$

Choose the corresponding complexity for both (a) and (b) and briefly justify your answer.

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Computation of the out-degree of a vertex $v$	$\Theta(1)$	$\Theta(\deg^+(v))$	$\Theta(\deg^-(v))$	$\Theta(n)$	$\Theta(m)$
Computation of the in-degree of a vertex $v$	$\Theta(1)$	$\Theta(\deg^+(v))$	$\Theta(\deg^-(v))$	$\Theta(n)$	$\Theta(m)$
Check for existence of an edge between two given vertices $u, v$	$\Theta(1)$	$\Theta(\deg^+(v))$	$\Theta(\deg^-(v))$	$\Theta(n)$	$\Theta(m)$
Check for existence of a 2-cycle starting at vertex $v$	$\Theta(n)$	$\Theta(\deg^+(v)\Delta^+)$	$\Theta(\deg^-(v)n)$	$\Theta(n^2)$	$\Theta(nm)$

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A 2-cycle is a cycle of length 2 containing two vertices  $u$  and  $v$  such that  $(u, v) \in E$  and  $(v, u) \in E$ .

### Task 21

Two nodes  $u$  and  $v$  in a graph are called *connected* if in the graph (1) there is a path from  $u$  to  $v$  and (2) a path from  $v$  to  $u$ . In undirected graphs, statements (1) and (2) are equivalent.

Develop an algorithm with runtime  $O(|V| + |E|)$  that assigns an (integer) label  $v.\ell$  to each vertex  $v$  of a given undirected graph  $G = (V, E)$  such that  $v.\ell = u.\ell$  for every pair of vertices  $u$  and  $v$  if and only if  $u$  and  $v$  are connected.